

# Designing pricing strategies for coordination of networked distributed energy resources

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**Abstract:** We study the problem of aggregator's mechanism design for controlling the amount of active, or reactive, power provided, or consumed, by a group of distributed energy resources (DERs). The aggregator interacts with the wholesale electricity market and through some market-clearing mechanism is incentivized to provide (or consume) a certain amount of active (or reactive) power over some period of time, for which it will be compensated. The objective is for the aggregator to design a pricing strategy for incentivizing DERs to modify their active (or reactive) power consumptions (or productions) so that they collectively provide the amount that the aggregator has agreed to provide. The aggregator and DERs' strategic decision-making process can be cast as a Stackelberg game, in which aggregator acts as the leader and the DERs are the followers. In previous work [Gharesifard et al., 2013b,a], we have introduced a framework in which each DER uses the pricing information provided by the aggregator and some estimate of the average energy that neighboring DERs can provide to compute a Nash equilibrium solution in a distributed manner. Here, we focus on the interplay between the aggregator's decision-making process and the DERs' decision-making process. In particular, we propose a simple feedback-based privacy-preserving pricing control strategy that allows the aggregator to coordinate the DERs so that they collectively provide the amount of active (or reactive) power agreed upon, provided that there is enough capacity available among the DERs. We provide a formal analysis of the stability of the resulting closed-loop system. We also discuss the shortcomings of the proposed pricing strategy, and propose some avenues of future work. We illustrate the proposed strategy via numerical simulations.

Keywords: Power systems, distributed energy resources, energy market, distributed control, game theory.

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## 1. INTRODUCTION

Power distribution networks are undergoing radical transformation in structure and functionality. These transformations are enabled by the increased reliance on advanced communications and controls, as well as by the increased penetration of renewable-based electricity generation resources (e.g., solar photovoltaics (PV) installations), controllable loads (e.g., thermostatically-controlled loads (TCLs)), and storage-capable loads (e.g., plug-in electric vehicles (PEVs)). These generation resources and loads are commonly referred to as distributed energy resources (DERs), and, if properly controlled, they can be utilized to provide ancillary services. For example, PEVs and TCLs can be utilized to provide frequency regulation services [Guille and Gross, 2009, Callaway and Hiskens, 2012]. However, in order to enable the added functionality that DERs may provide, it is necessary to develop appropriate control mechanisms. In this paper, we address this problem and propose a framework for controlling the power provided/consumed by DERs, and perhaps also the reactive power if the objective is to regulate voltage.

\* This work was supported in part by a grant through the Information Trust Institute of the University of Illinois; and by NSF under grant ECCS-CPS-1135598.

Focusing on controllable loads, their control is currently achieved through demand response programs in which participants sign a contract with an aggregating entity—the demand response provider—so that their electrical energy consumption can be curtailed by the aggregator in response to market prices or in order to ensure reliable operation of the system, in exchange for lower electricity prices. In this work, we also consider an aggregating entity that will interact with the wholesale electricity market and, through pricing, will incentivize DERs to provide/consume active (or reactive) power in exchange for monetary benefits. As an example, a household with a PV system (with a reactive power capable power electronics grid interface) and a TCL might choose to offer these two resources to an aggregator so that the PV system is utilized to provide reactive power for voltage control, and the TCL is utilized to provide frequency regulation. In this sense, the interplay between the DERs and the aggregator can be modeled as a Stackelberg *game*, where the aggregator acts as the leader and the DERs are the followers.

### 1.1 Literature Review

Game-theoretic models have been used recently for studying energy markets (see, e.g., [Fan, 2012, Kiani and An-

naswamy, 2012, Saad et al., 2012]). The game-theoretic parts of our work are related to noncooperative resource allocation problems in the literature, for example [Johari and Tsitsiklis, 2006, Maheswaran and Başar, 2006, Kelly et al., 1998]), where under appropriate concavity assumptions, the existence of Nash equilibrium in pure strategies is guaranteed using the results in [Rosen, 1965]. In the context of our work, the process of DER power consumption is related to the charging process in PEVs that recently appeared in [Tushar et al., 2012], and [Ma et al., 2012]; however, in our setting, we deal with a scenario in which the DERs are individual decision makers and compute the Nash equilibrium using the information available from their neighboring DERs (along with the price set by the aggregator). Additionally, in our setting, we include more general pricing strategies and allow for analysis of scenarios in which DERs are capable of both providing and consuming power.

### 1.2 Statement of Contributions

The contributions of this paper are as follows. First, building on the framework introduced in our recent work [Gharesifard et al., 2013b,a], we consider a two-layer decision-making structure. In the first layer—the aggregator layer—there is a set of aggregators that interact with the wholesale electricity market, and through some market-clearing mechanism, will enter a contract to provide a certain amount of active (or reactive) power over some predetermined period of time, for which they will be compensated. In the second layer—the retail market layer—each aggregator offers prices for the active (or reactive) power that DERs may provide or consume—negative if provided and positive if consumed. Then, the objective is for the aggregator to incentivize the DERs via pricing to provide or consume active (or reactive) power so that they collectively provide an amount equal to that the aggregator has agreed to provide. We then focus on the retail market layer; specifically, on the interplay between an aggregator and a set of DERs.

We assume that each DER is a price-anticipating decision-maker (i.e., it is aware of the aggregator pricing strategy), and in order to make a decision, it exchanges some information with neighboring DERs, with the objective of estimating the average amount of resource (active or reactive power) collectively available in its immediate neighborhood. Then, each DER uses this estimate, together with the aggregator’s pricing strategy information, to decide on the amount of active (or reactive) power that it will sell or buy; this strategic decision-making process is cast as a game which, under appropriate assumptions is guaranteed to have a unique Nash equilibrium solution. When the collective active (or reactive) power consumed (or produced) by the DERs does not match the amount that the aggregator has been asked to provide, the aggregator will adjust its pricing strategy so as to steer the DERs to collectively provide the amount that the aggregator has contracted in the wholesale electricity market. In particular, we provide a feedback control mechanism that allows the aggregator to change its pricing strategy to meet its objective, i.e., to coordinate the DERs so that they collectively provide the amount of active (or reactive) power that the aggregator has agreed to provide. The upshot of this feedback control mechanism described above is that it preserves privacy, in

the sense that it does not rely on the utility functions of the individual DERs.

## 2. MATHEMATICAL PRELIMINARIES

We start with some notational conventions. Let  $\mathbb{R}$ ,  $\mathbb{R}_{\geq 0}$ ,  $\mathbb{Z}$ , and  $\mathbb{Z}_{\geq 1}$  denote the set of real, nonnegative real, integer, and positive integer numbers, respectively. We denote by  $\mathcal{B}(X)$  the set of bounded real-valued functions on a set  $X \subset \mathbb{R}^d$ ,  $d \in \mathbb{Z}_{\geq 1}$ ; we use  $\mathcal{B}^0(X)$  when the functions are, additionally, continuous. We use the short-hand notation  $\mathbf{1}_d = (1, \dots, 1)^T \in \mathbb{R}^d$  and  $\mathbf{0}_d = (0, \dots, 0)^T \in \mathbb{R}^d$ . We also denote by  $e_i \in \mathbb{R}^d$  the unit vector with 1 in its  $i$ th coordinate and zeros elsewhere.

### 2.1 Graph Theory

Next, we introduce some basic notions from graph theory [Biggs, 1994]. A *graph*  $\mathcal{G} = (\mathcal{V}, \mathcal{E})$  consists of a *vertex* set  $\mathcal{V}$  and an *edge* set  $\mathcal{E} \subseteq [\mathcal{V}]^2 := \{(v, w) \mid v, w \in \mathcal{V}\}$ . The *neighbors* of  $v \in \mathcal{V}$  belong to the set  $\mathcal{N}_v := \{w \in \mathcal{V} \mid (v, w) \in \mathcal{E}\}$  and the *degree* of  $v$  is  $d_v := |\mathcal{N}_v|$ . A *path* in  $\mathcal{G}$  is a sequence of edges of the form  $(v_1, v_2), (v_2, v_3), \dots, (v_{k-1}, v_k)$ . The graph  $\mathcal{G}$  is *connected* if there is a path between any pair of vertices. Henceforth, when not explicitly stated, we fix an ordering on the vertex set  $\mathcal{V}$  so that we may write  $\mathcal{V} = \{v_1, \dots, v_n\}$ , where  $n$  is the order of  $\mathcal{G}$ . The *adjacency matrix* of  $\mathcal{G}$  is the  $n \times n$  matrix  $\mathbf{A}$  defined as  $A_{i,j} = 1$  if  $\{v_i, v_j\} \in \mathcal{E}$  and  $A_{i,j} = 0$  otherwise. The weighted degree  $v_i$ ,  $i \in \{1, \dots, n\}$  is  $d(v_i) = \sum_{j=1}^n a_{ij}$ . The *weighted degree matrix*  $\mathbf{D}$  is the diagonal matrix defined by  $(\mathbf{D})_{ii} = d(i)$ , for all  $i \in \{1, \dots, n\}$ . The *Laplacian* is  $\mathbf{L} = \mathbf{D} - \mathbf{A}$ . For an (undirected) graph,  $\mathbf{L}\mathbf{1}_n = \mathbf{1}_n^T \mathbf{L} = \mathbf{0}$ , and  $\mathbf{L} = \mathbf{L}^T$  and is positive semidefinite. When  $\mathcal{G}$  is connected, the zero eigenvalue is simple.

### 2.2 Game Theory

We recall the class of concave games in the absence of shared constraints, see [Rosen, 1965, Başar and Olsder, 1999]. A *concave game* (with unshared constraints) is a triplet  $\mathbf{G} = (V, S, \{f_i\}_{i=1}^n)$ , where

- $V$  is a group of  $n \in \mathbb{Z}_{\geq 1}$  players,
- $S = S_1 \times S_2 \times \dots \times S_n$  is the product strategy set, where  $S_i \subset \mathbb{R}^{d_i}$  with  $d_i \in \mathbb{Z}_{\geq 1}$ , is nonempty, convex and compact, and
- $f_i : S \rightarrow \mathbb{R}$ , the payoff for player  $i \in \{1, \dots, n\}$ , is a locally Lipschitz concave mapping.

A point  $x^* \in S$  is called a *Nash equilibrium* of  $G$  if and only if, for all  $i \in V$ ,

$$f_i(x^*) = \max_{y_i} \{f_i(x_1^*, \dots, x_{i-1}^*, y_i, x_{i+1}^*, \dots, x_n^*) \mid y_i \in S_i\}.$$

In other words, when the game is at  $x^*$ , no player can improve its payoff by unilaterally deviating from this point. The existence of a (pure) Nash equilibrium for this class of games is guaranteed [Rosen, 1965]. A uniqueness result can also be obtained under the so-called diagonally strict concavity assumption, along with differentiability (see [Rosen, 1965, Theorem 4]), when one considers another suitable notion of equilibrium (the so-called normalized or variational equilibrium), see [Kulkarni and Shanbhag, 2012]. When the constraints are not shared, as it is the case in this paper, these notions of equilibria match, yielding an applicable uniqueness result.

### 3. PROBLEM STATEMENT

We first recall some of the notions that were introduced in [Gharesifard et al., 2013b,a]. Consider a set of aggregators, denoted by  $\{v_1^{\text{agg}}, \dots, v_N^{\text{agg}}\}$ ,  $N \in \mathbb{Z}_{\geq 1}$ , who, through some market-clearing mechanism, are asked to provide a certain amount of resource (active or reactive power) over some predetermined period of time. We assume that each aggregator has a backup storage device that can be used to provide (or consume) the difference. Alternatively, one can assume that the aggregator contracts some insurance with a third party that will provide the difference. Each aggregator is responsible for controlling the amount of resource provided or consumed by a group of DERs via some pricing strategy. The collection of all DERs builds a new layer, which we term the *retail market layer*. The notions above are illustrated in Fig. 1, with the decision-making process of one of the aggregators as the focus of the paper; the decision-making of DERs in the retail market layer was discussed in [Gharesifard et al., 2013b,a], as we recall next.

#### DERs' decision-making process

We denote by  $V = \{v_1, \dots, v_n\}$ ,  $n \in \mathbb{Z}_{\geq 1}$ , the set of DERs interacting with the aggregator, where the available amount of resource of each  $v_i$ ,  $i \in \{1, \dots, n\}$ , at time  $t \in \mathbb{R}_{\geq 0}$  is denoted by  $x_i(t) \in [0, 1]$ . Each DER is a decision maker and can freely choose to participate after receiving a request from its aggregator. The DERs can remain idle, or provide/absorb active (or reactive) power. The decision that each DER is faced with, among other things, depends on its own utility function, along with the pricing strategy designed by the aggregator. The DERs considered in [Gharesifard et al., 2013b,a], and also here, are price anticipating, in the sense that they are aware of the fact that the pricing is designed by the aggregator with respect to the average resource available. We also assume that in order for each DER to make its decision, it can collect information from neighboring DERs with which it can exchange information. Let us formalize these statements.

Without loss of generality, we assume that the DER is willing to participate in the resource reallocation process by setting the amount of resource that it is providing in the range  $[0, 1]$ . The functions  $P_c : [0, 1] \rightarrow \mathbb{R}_{\geq 0}$  and  $P_p : [0, 1] \rightarrow \mathbb{R}_{\geq 0}$  given by  $P_c(\bar{x}(t))$  and  $P_p(\bar{x}(t))$ ,  $\bar{x}(t) = \frac{1}{n} \sum_{i=1}^n x_i(t)$ , denote, respectively, the price per unit of resource that the DERs pay when consuming and receive when providing. Additionally, it is reasonable to assume that the DERs are making their decisions on the amount of resource they will provide or absorb for a certain a priori fixed period of time that they are aware of. We also need to define a set of utility functions  $U_i : [0, 1] \rightarrow \mathbb{R}_{\geq 0}$ , with values  $U_i(x_i)$ , for all  $i \in \{1, \dots, n\}$ . Although our treatment allows for more general classes of utility functions, it is assumed that each  $U_i$  is an increasing function of the available resource. Let us next describe the decision-making process that each DER is faced with. Similar to other scenarios of resource allocation problems (see, e.g., [Johari and Tsitsiklis, 2006]), each DER wishes to maximize a payoff function  $f_i : X \times X_{\text{agg}} \rightarrow \mathbb{R}$ , where  $X = [0, 1]^n$  and  $X_{\text{agg}} = \mathcal{B}^0([0, 1]) \times \mathcal{B}^0([0, 1])$ , given by

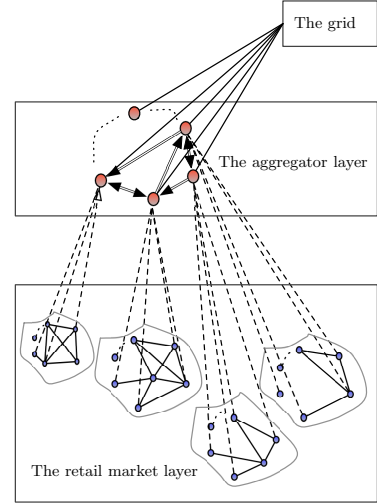


Fig. 1. Networks of DERs and their corresponding aggregators; in each network, each DER can communicate with some neighbors. DERs can also communicate with their corresponding aggregator.

$$f_i(x_i, x_{-i}, P_c, P_p) = \begin{cases} U_i(x_i) - (x_i - x_i^0)P_c(\bar{x}), & x_i > x_i^0, \\ U_i(x_i) - (x_i - x_i^0)P_p(\bar{x}), & x_i \leq x_i^0, \end{cases} \quad (1)$$

where  $(x_i^0, x_{-i}^0) \in X$  denotes the initial resource profile of all DERs.

#### Aggregator's objective

Let  $\mathcal{X} \in \mathbb{R}$  be the amount of resource (active or reactive power) that the aggregator has contracted to provide/consume over some predetermined period of time; when  $\mathcal{X} \in \mathbb{R}_{<0}$ , the aggregator needs to encourage the DERs to provide active (or reactive) power. Conversely, when  $\mathcal{X} \in \mathbb{R}_{>0}$ , the aggregator needs to encourage the DERs to consume active (or reactive) power. The objective of the aggregator is to ensure that the DERs collectively provide  $\mathcal{X} \in \mathbb{R}$  units of resource; thus it wishes to maximize the function  $f_{\text{agg}} : X \times X_{\text{agg}} \rightarrow \mathbb{R}_{\leq 0}$  given by

$$f_{\text{agg}}(x, P_c, P_p) = -|\mathcal{X} - \sum_{i=1}^n \alpha_i(x_i - x_i^0)|,$$

where  $\alpha_i \in \mathbb{R}_{>0}$ , for all  $i \in \{1, \dots, n\}$ , over the choices of  $P_c$  and  $P_p$ .

Based on the description given above, the aggregator and the DERs define a game, called the *retail market game*,

$\mathbf{G}_{\text{DERs-AGG}} = (V \cup \{v^{\text{agg}}\}, \mathbb{R}^n \times X_{\text{agg}}, f_1 \times \dots \times f_n \times f_{\text{agg}})$ , where players wish to maximize their objective functions. For most part, we think of  $\mathbf{G}_{\text{DERs-AGG}}$  as a two-stage Stackelberg game, where the aggregator sets the pricing strategies first and the DERs act next.

We next raise some questions, where (a-c) were previously introduced in [Gharesifard et al., 2013b,a].

- (a) (*Existence of equilibria*): given the pricing strategies of the aggregator  $P_c, P_p \in \mathcal{B}^0([0, 1])$ , does there exist a Nash equilibrium solution to the retail market game? If so, is the equilibrium unique?



$$\begin{aligned}
\dot{\mathbf{x}}(t) &= F(\mathbf{x}(t), \mathbf{z}(t)) + G\xi(t), \\
\dot{\mathbf{z}}(t) &= \mathbf{L}\mathbf{x}(t), \\
\dot{\xi}(t) &= -\left(\sum_{i=1}^n (x_i(t) - x_i^0) - \mathcal{X}\right), \quad (5)
\end{aligned}$$

asymptotically converges to  $(\mathbf{x}^*, \mathbf{z}^*)$ .

**Proof.** For simplicity of notation, let us introduce

$$\eta(t) = -\left(\sum_{i=1}^n (x_i(t) - x_i^0) - \mathcal{X}\right),$$

where  $(x_1, \dots, x_n) \in \mathbb{R}^n$  is the resource profile of the DERs. For an initial condition  $(\mathbf{x}^0, \mathbf{z}^0, \xi_0)$ , consider the candidate Lyapunov function  $V_1 : X^n \times \mathbb{R}^n \times \mathbb{R} \rightarrow \mathbb{R}_{\geq 0}$  given by

$$V_1(\mathbf{x}, \mathbf{z}, \xi) = V_0(\mathbf{x}, \mathbf{z}) + \frac{1}{2}(\xi - k_d)^2,$$

where  $V_0$  is a quadratic control Lyapunov function for (2) given by

$$V_0(\mathbf{x}, \mathbf{z}) = \frac{1}{2}(\mathbf{x} - \mathbf{x}^*)^T(\mathbf{x} - \mathbf{x}^*) + \frac{1}{2}(\mathbf{z} - \mathbf{z}^*)^T(\mathbf{z} - \mathbf{z}^*),$$

with the stabilizing feedback  $P(\bar{\mathbf{x}}(t)) = a\bar{\mathbf{x}}(t) + k_d$ . The function  $V_1$  is clearly a smooth mapping. Let us examine its set-valued Lie derivative.

$$\begin{aligned}
\mathcal{L}_\Psi V_1(\mathbf{x}, \mathbf{z}, \xi) &= \frac{\partial V_0^T}{\partial \mathbf{x}} (F(\mathbf{x}(t), \mathbf{z}(t)) + G\xi(t)) + \\
&+ \frac{\partial V_0^T}{\partial \mathbf{z}} \mathbf{L}\mathbf{x}(t) + (\xi(t) - k_d)\eta(t) \\
&= \frac{\partial V_0^T}{\partial \mathbf{x}} (F(\mathbf{x}(t), \mathbf{z}(t)) + Gk_d) + \frac{\partial V_0^T}{\partial \mathbf{z}} \mathbf{L}\mathbf{x}(t) \\
&+ (\xi(t) - k_d)\eta(t) + \frac{\partial V_0^T}{\partial \mathbf{x}} G(\xi(t) - k_d) \\
&= \mathcal{L}_\Phi V_0(\mathbf{x}, \mathbf{z}) + (\xi(t) - k_d)(\eta(t) + \frac{\partial V_0^T}{\partial \mathbf{x}} G).
\end{aligned}$$

By definition of  $G$ , we have that  $\eta(t) + \frac{\partial V_0^T}{\partial \mathbf{x}} G = 0$ . Given that  $\mathcal{L}_\Phi V_0(\mathbf{x}, \mathbf{z}) \leq 0$ , we conclude that

$$\mathcal{L}_\Psi V_1(\mathbf{x}, \mathbf{z}, \xi) \leq 0,$$

and hence (5) is stable and has bounded trajectories. The application of the LaSalle invariance principle leads to the conclusion that, from a given initial condition, the evolution of (5) approaches a set  $M$  of the form  $V_1^{-1}(c) \cap S$ ,  $c \in \mathbb{R}$ , where  $S$  is the largest invariant set contained in  $\{(\mathbf{x}, \mathbf{z}, \xi) \in V_1^{-1}(\leq V_1(\mathbf{x}^0, \mathbf{z}^0, \xi_0)) \mid \mathcal{L}_\Psi V_1(\mathbf{x}, \mathbf{z}, \xi) = 0\}$ .

Note, however, that when  $\mathcal{L}_\Psi V_1(\mathbf{x}, \mathbf{z}, \xi) = 0$ , we have that  $\mathcal{L}_\Phi V_0(\mathbf{x}, \mathbf{z}) = 0$ , which implies that  $(\mathbf{x}, \mathbf{z}) = (\mathbf{x}^*, \mathbf{z}^*)$ . ■

**Remark 5.2. (Implementation of the continuous-time controller):** Although we have designed the pricing inputs as a continuous-time controller, it can potentially be implemented in a discrete-time manner; this is a reasonable assumption as the change in pricing happens much slower than the Nash seeking dynamics in the retail market. A real-time scheduling can also be achieved using an event-triggered implementation [Tabuada, 2007]. Finally, it is worth mentioning that the procedure proposed is robust, a property of the saddle-point dynamical systems, and can regulate slow time-varying requests. •

**Remark 5.3. (Optimal regulating pricing strategy):** The shortcoming of the pricing strategy provided here is that it does not take into account the revenue that

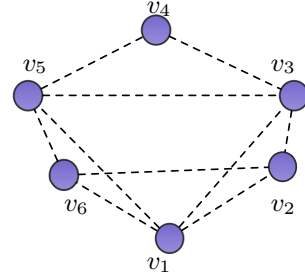


Fig. 2. The network of DERs of the simulations is shown. The adjacency matrix associated with this network has 1 in the entries corresponding to edges and zeros elsewhere.

	DER 1	DER 2	DER 3	DER 4	DER 5	DER 6
$u^1$	0.4021	0.7126	0.5502	0.6002	0.8216	0.7984
$u^2$	0.1000	0.1000	0.1000	0.1000	0.1162	0.1000

Table 1. The utility properties of each DER in Fig. 2.

the aggregator obtains. This is because the aggregator typically receives a fixed price through the contract in the wholesale market, but this price might be lower than the one that the aggregator needs to pay the DERs. Problem statement (d) takes this into account, by modeling the aggregators' decision-making process as the leader of a Stackelberg game. The difficulty, however, is in designing an implementable optimal solution for the aggregator which does not rely on the private information about the utility functions of the DERs. The approach developed in [Shen and Başar, 2007] provides promising avenues for addressing this. •

## 6. SIMULATIONS

Consider a group of DERs connected to an aggregator; for illustration purposes, we have only selected six DERs  $\{v_1, \dots, v_6\}$ . The DERs are price anticipating and can obtain information from each other via a communication network shown in Fig. 2.

Each DER's utility function  $U_i : [0, 1] \rightarrow \mathbb{R}_{\geq 0}$ ,  $i \in \{1, \dots, 6\}$ , is given by

$$U_i(x) = u_i^1 \log(1+x) + u_i^2 x,$$

where  $U_i$  is normalized so that  $u_i^1, u_i^2 \in (0, 1]$ . The values of these parameters for each DER are given in Table 1. Note that  $U_i$  is increasing and strictly concave, and thus satisfies Assumption 4.1(i). The pricing strategy for the aggregator is given by (3) and is thus linear. The parameter  $a$  is chosen to be 0.1 here. It can be shown that with the given parameters, the diagonally strict concavity assumption holds.

Fig. 3 demonstrates the evolutions of the dynamics (5) from the initial condition  $x_0 = (0.2, 0.3, 0.3, 0.4, 0.1, 0.2)^T$ . In the first scenario, see (a) and (b), the aggregator is requested to encourage the DERs to consume resource such that 50% of their total capacity is full. The pricing input of the aggregator is shown in Fig. 4(a). In contrast, in Fig. 3(c) and (d), the aggregator encourages the DERs to provide one-third of their available resource back to the grid. Fig. 4(b) shows the pricing input of the aggregator for this case. The pricing control inputs of this example can be implemented in discrete-time, where the aggregator implements changes in its pricing scheme in a longer period of time, using its information about the average available resource within the DERs.



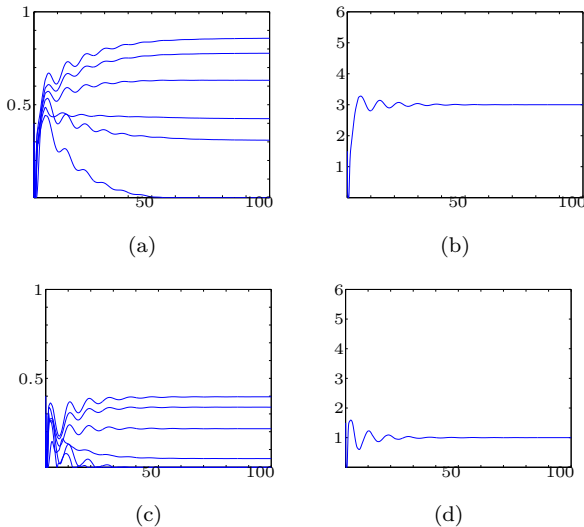


Fig. 3. (a) and (b) correspond to the scenario where the grid has requested that the aggregator encourages the DERs to consume resource such that 50% of their total capacity is used. The DERs' initial charge is  $x_0 = (0.2, 0.3, 0.3, 0.4, 0.1, 0.2)^T$ , i.e., 25% of their total capacity is used. Panel (a) shows each DERs' consumption, and Panel (b) shows their total capacity, regulated by the aggregator. In contrast, (c) and (d) correspond to the scenario where the grid has requested that the aggregator encourages the DERs to return one-third of their available resource (active or reactive power) to the grid.

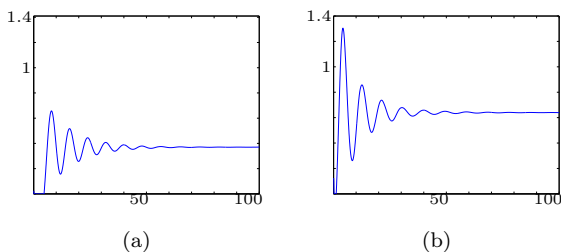


Fig. 4. (a) and (b) show the pricing control input of the aggregator, asymptotically converging to 0.3704 and 0.6396, respectively.

## 7. CONCLUSIONS AND FUTURE WORK

We have studied the problem of incentivizing a group of DERs to provide/consume active or reactive power in exchange for monetary benefits. Building on the framework previously introduced in [Gharesifard et al., 2013b,a], the DERs considered are price anticipating and only have access to local information, but can still compute an equilibrium solution once the pricing strategy is fixed.

We have proposed a pricing strategy that allows the aggregator to regulate the consumption (or provision) of active (or reactive) power by a group of DERs. The upshot of the pricing strategy proposed is that it does not rely on the private information about the utilities of the DERs.

We have also discussed a possible shortcoming of this strategy regarding the revenue of the aggregator, and have proposed a framework for addressing it. It is an interesting avenue of research to study this problem for classes of nonlinear pricing strategies as in [Shen and Başar, 2007], where the aggregator also uses the average-dependent pricing terms for regulation.

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